

A Survey of Algorithms

Alternatives to Buchberger's Algorithm

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Significance of Grobner Bases

- 1 Hilbert's 10th Problem
- 2 Find an algorithm to determine whether a given polynomial Diophantine equation with integer coefficients has an integer solution
- 3 (1970) Deemed **Impossible** by Matiyasevich's Theorem (MRDP Theorem)
- 4 Grobner Bases: Solve Problems that are considered computationally hard

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Introduction

- 1 The **input** is a finite set of polynomials, and **output** is a finite Grobner basis.
- 2 Buchberger's Algorithm requires the use of **S-polynomial** and **Division Algorithm**
- 3 Recall

$$S(f_1, f_2) = \frac{M}{\text{LT}(f_1)} f_1 - \frac{M}{\text{LT}(f_2)} f_2$$

where $M = \text{LCM}((\text{LT}(f_1), \text{LT}(f_2)))$.

Problems With Buchberger's Algorithm

- 1 Simplicity, Efficiency, Memory Usage
- 2 Many "useless" S -polynomial computation (several divisions that reduce to 0)
- 3 Buchberger's product and chain criterion to reduce complexity (1979, 1985)

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Important Features of Other Algorithms

- 1 **Product Criterion:** If $\text{LCM}(\text{LT}(f), \text{LT}(g)) = \text{LT}(f)\text{LT}(g)$, then the pair (f, g) can be removed.
- 2 **Chain Criterion:** A pair (f, g) can be removed if there is some h such that $\text{LT}(h) \mid \text{LCM}(\text{LT}(f), \text{LT}(g))$, and both pairs (f, h) and (h, g) have been removed before.
- 3 For a fixed polynomial r -tuples $\mathbf{f} = (f_1, \dots, f_r) \in P^r$, $(g_1, \dots, g_r) \in P^r$ is called a **syzygy** wrt \mathbf{f} if $\sum g_i f_i = 0$.
- 4 Sparse matrices, "signature," and rewriting

Historical Progress

- 1 Moller, Mora, Traverso present an algorithm that uses full module of syzygies, but inefficient (1992)
- 2 Grobner Walk: Conversion between Grobner basis for different monomial orders
- 3 Faugere's F4 and F5 Algorithm
- 4 F4: Normal forms computed and makes use of sparse matrices (1999). Easy to understand, efficient, but memory usage grows quickly
- 5 F5 algorithm **detects all** useless S-polynomial reductions (2002) via signatures and rewriting rules. Efficient but difficult to understand
- 6 Last 15 years, G2V, GVW, and other variants

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Computer Programs

- 1 FGb, Maple
- 2 CoCoA, Macaulay2, Magma, Singular, Sage
- 3 Mathematica: Buchberger and Groebner walk
- 4 Maple: fgb, maplef4 (F4 algorithm), buchberger, fglm (Faugere, Gianni, Lazard, Mora), Groebner walk

Maple Demonstration

Groebner base from tdeg order to plex order

Asked 2 years, 6 months ago Active 1 year, 4 months ago Viewed 90 times

▲ I try to solve an equation by using Groebner bases. When I use Maple to find its Groebner basis with *plex order*, Maple take too long to calculate and the proceed does not terminate. Thus, I try to find with *tdeg order* and the proceed takes seconds.

▼ Now, I have the Groebner basis with *tdeg order*. I want to use it to find the Groebner basis with *plex order*. Can I do it in Maple?

2

1



maple groebner-basis

Mathematica Demonstration

1 Consider

$$\{xy^4 + yz^4 - 2x^2y - 3, y^4 + xy^2z + x^2 - 2xy + y^2 + z^2, -x^3y^2 + xyz^3 + y^4 + xy^2z - 2xy\}.$$

`GroebnerBasis[polys, {x, y, z}, Method -> "GroebnerWalk"]; // Timing`

2 {0.1875, Null}

`TimeConstrained[GroebnerBasis[polys, {x, y, z}, Method -> "Buchberger"];, 60]`

3 \$Aborted

Maple Results

1 $\{x^2 - 2xz + 5, xy^2 + yz^3, 3y^2 - 8z^3\}$

`> Basis(F, lexdeg([x], [y,z]), method=walk);`

-> Groebner Walk

total time: 0.003 sec

$$\left[8z^3 - 3y^2, \right.$$

$$\left. 9y^4 + 48y^3z + 320y^2, 8xy^2 + 3y^3, x^2 - 2xz + 5 \right]$$

2

`> Basis(F, grlex(x,y,z));`

-> F4 algorithm

total time: 0.005 sec

$$\left[x^2 \right.$$

$$\left. -2xz + 5, 8z^3 - 3y^2, 8xy^2 + 3y^3, 9y^4 + 48y^3z + 320y^2 \right]$$

3

total time: 0.013 sec

$$\left[x^2 \right.$$

$$\left. -2xz + 5, 8z^3 - 3y^2, 8xy^2 + 3y^3, 9y^4 + 48y^3z + 320y^2 \right]$$

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Applications

- 1 Faugere's F5 solved the first Hidden Field Equation (HFE) Cryptosystem Challenge (80 polynomial equations with 80 unknowns)
- 2 Cyclic 10 Problem solved by F5
- 3 Cryptography, robotics, celestial mechanics, signal theory, error correcting codes
- 4 Other related algorithms: Knuth-Bendix Completion, Robinson's resolution in automated theorem proving

The End

- Thank You!
- Questions?